

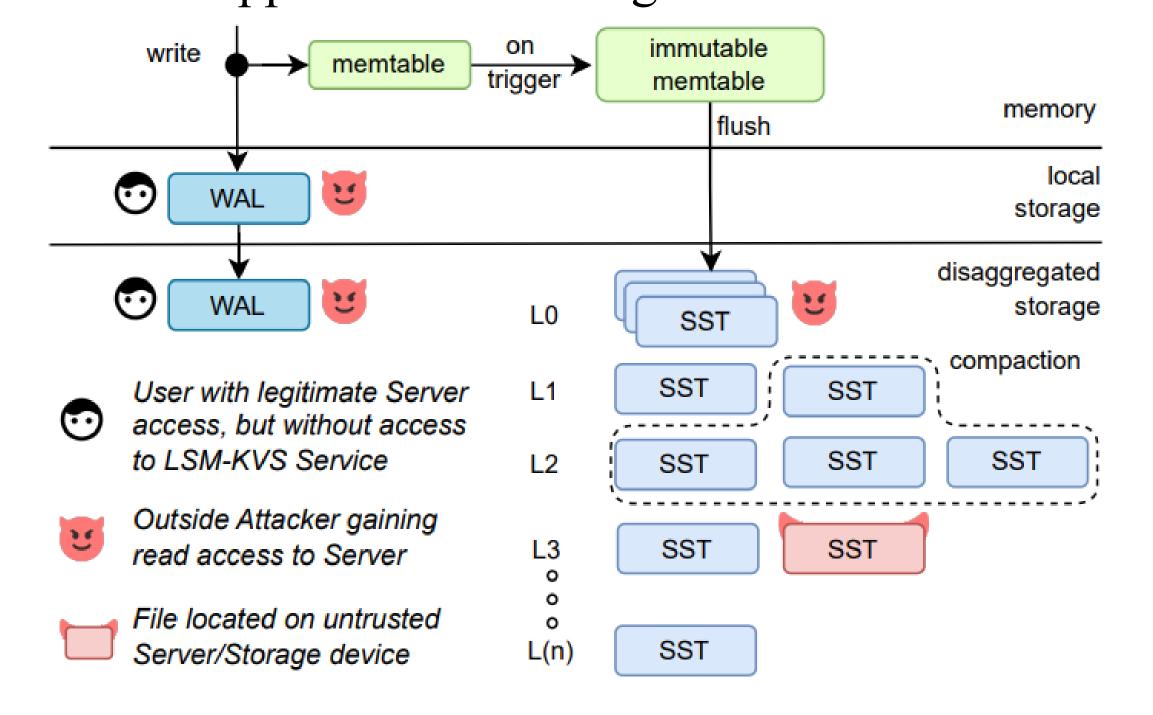
# SHIELD: Encrypting Persistent Data of LSM-KVS from Monolithic to Disaggregated Storage



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# THREAT MODEL

Plaintext LSM-KVS files on multiple servers with untrusted applications and large attack surface.



## MOTIVATION

- LSM-KVS in disaggregation has components (compaction, storage, WAL) across servers. Solution must be flexible to disaggregated setups.
- Data Encryption Key (DEK) practices (Unique DEK per file and DEKrotation) necessary for robust protection in disaggregation.
- SOTA presents high overhead (350-3,750%) from using hardware-based TEEs. It's focus on in-memory protection misses bottlenecks (WAL-write), while using a single DEK.

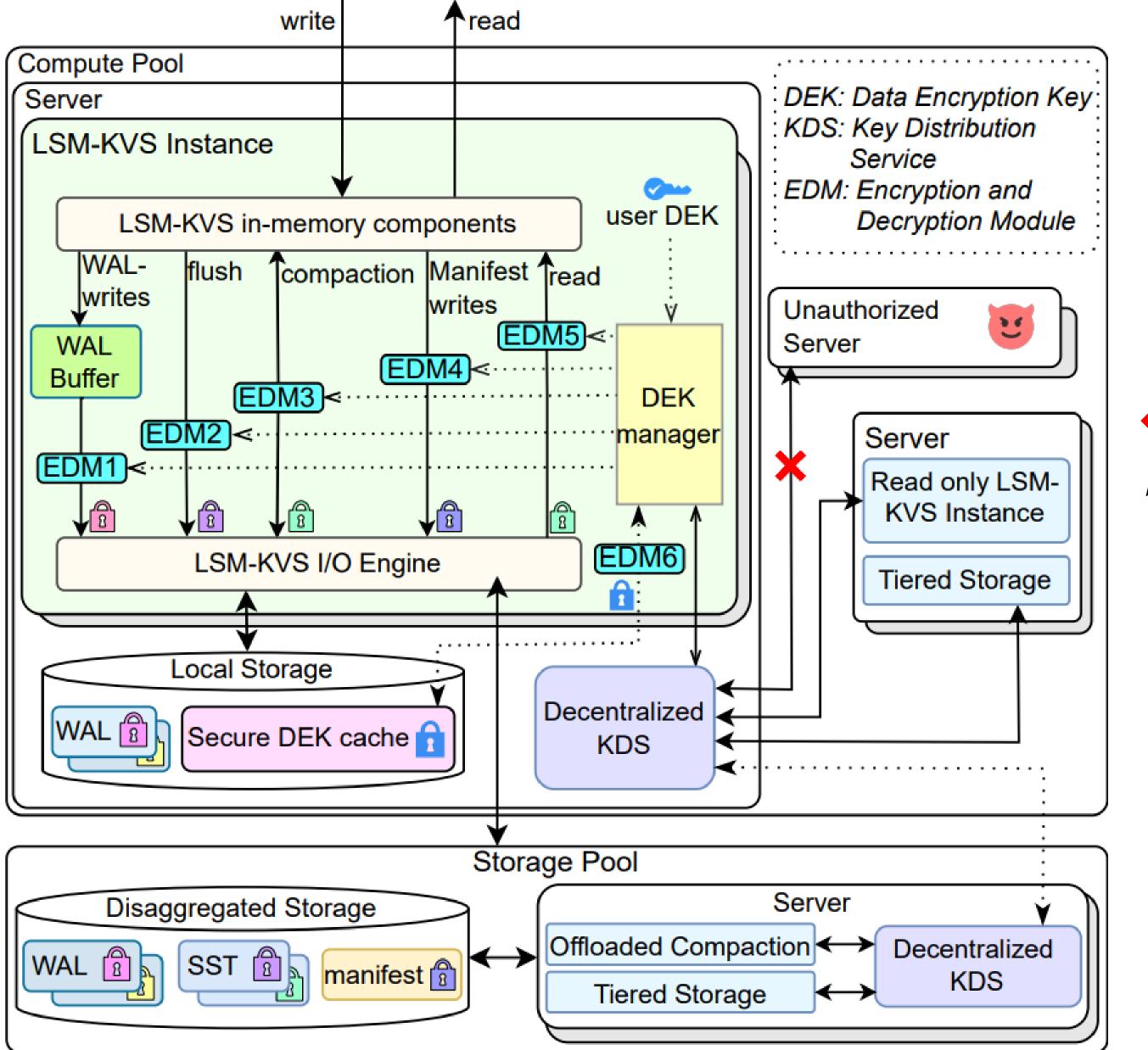
### CHALLENGES

- 1. How and where to embed DEK-practices in LSM-KVS?
- How to mitigate the WAL-Write Encryption Bottleneck?
- How to co-ordinate DEK/LSM-file relationship for flexible setups?

# ENCFS (DISTRIBUTED) & SHIELD (DISAGGREGATED STORAGE)

### EncFS (For Monolith and Distributed):

- Intercept IO to/from FS to enable transparent encryption in LSM-KVS.
- Simple solution but DEK leak can compromise all data.
- For servers completely in your control. Can be killed on-demand.



Unencrypted RocksDB + WAL-But

#### ♠ read LSM-KVS Instance LSM-KVS LSM-KVS Instance Instance LSM-KVS in-memory components (IMC) IMC IMC Manifest-WAL-Compaction writes IOE 0 0 0 LSM-KVS I/O Engine (IOE) TEM **←** TEM Transparent Encryption Module (TEM) user-provided data encryption key (DEK) Encrypted LSM-KVS Persistent files Other WAL 宜 processes SST 🛐 conf 宜

### SHIELD (For Disaggregated):

- SST encrypted after LSM block chunking. And integration into Compaction path for DEK-rotation.
- Size-configurable WAL-buffer to mitigate WAL-bottleneck
- Decentralized key distribution service (e.g., SSToolkit, Kerberos) for DEK provision the utilizes unique DEK identifiers.
- Local DEK-cache secured with user-selected password for other LSM instances to use, avoiding network round-trips.
- DEK-identifiers stored in LSM-KVS files' metadata. Authorized servers avoid lookup tables, requesting DEK from KDS for passive and flexible DEK sharing.

# EVALUATION RESULTS

#### 16-33% (EncFS) 0-1% **6-14**% 3-6% 9-36% (SHIELD) (50, 50)(50, 50)(99, 1)(write, read) ratio (write, read) ratio **Monolith Setup Disaggregated Storage Converging Trend for EncFS & SHIELD Compatible with different** → SHIELD Unencrypted RocksDB **Compaction Styles** Unencrypted RocksDB + WAL-Buf Throughput (kops/s) 00 05 Throughput (kops/s) Universal 200 500 Level 100 1000 No. of Writer Threads **FILLRANDOM** Value Size (Bytes) -X- Unencrypted RocksDB → EncFS - SHIELD Throughput (kops/s) Throughput (kops/s)

### DISCUSSION

- We use 128-bit AES in CTR mode for encryption. SHIELD is compatible with other encryption algorithms.
- The solution comes with 0-36% overhead. Encryption has a cost. SHIELD aims to avoid the penalty and reduce round trips.
- We promise updates for 3 RocksDB major revisions for the SHIELD codebase.

### FUTURE WORK

- What's the deal with TEEs? SOTA suffers high penalties, can TEE-penalty be avoided? Exploration needed.
- SHIELD focuses of DS. Disaggregated Memory with CXL or RDMA devices will be intriguing future work.



READRANDOM

Universal





Buffer Size (Bytes)

512 1024

2048

1024

Buffer Size (Bytes)

ASU-IDI

2048